

The Zonnon Object Model: A Structured Approach to Composability & Concurrency

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The Pascal Language Family

- Guiding Principle: „Make it as simple as possible but not simpler“

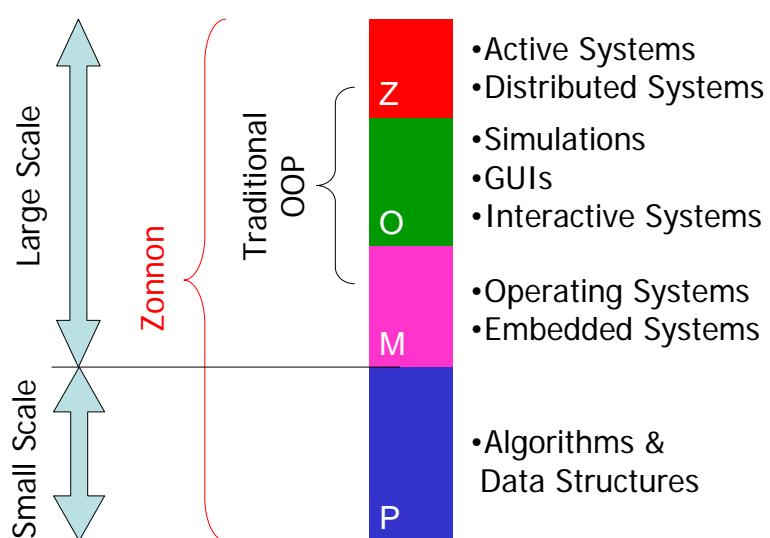
	Language	New Feature	Concept
1970	Pascal	Pointer	A&D
1980	Modula	Module	Systems
1990	Oberon	Type Extension	OOP
2005	Zonnon	Activity	Concurrency

- The Zonnon project has emerged from MS Project 7/7+ initiative
<http://zonnon.ethz.ch>

“With a new computer language,
one not only learns a new
vocabulary and grammar but one
opens oneself to a new world of
thought”

Niklaus Wirth

Spectrum of Programming



The Object Model

- Modular
- Compositional
- Active

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A Simple Interactive Module

- ```
module Weekday imports Zeller;
procedure { public } Compute;
 var d, m, y: integer;
begin
 read(d);
 while d > 0 do
 readln(m, y);
 writeln(Zeller.Weekday(
 y div 100, y mod 100, m - 2, d));
 read(d)
 end
end Compute;
end Weekday.
```

## A Simple Interactive Module

- ```
module Weekday imports Zeller;
procedure { public } Compute;
  var d, m, y: integer;
begin
  read(d);
  while d > 0 do
    readln(m, y);
    writeln(Zeller.Weekday(
      y div 100, y mod 100, m - 2, d));
    read(d)
  end
end Compute;
end Weekday.
```

no classes,
no objects,
no inheritance,
no virtual methods,
no overriding,
no static fields

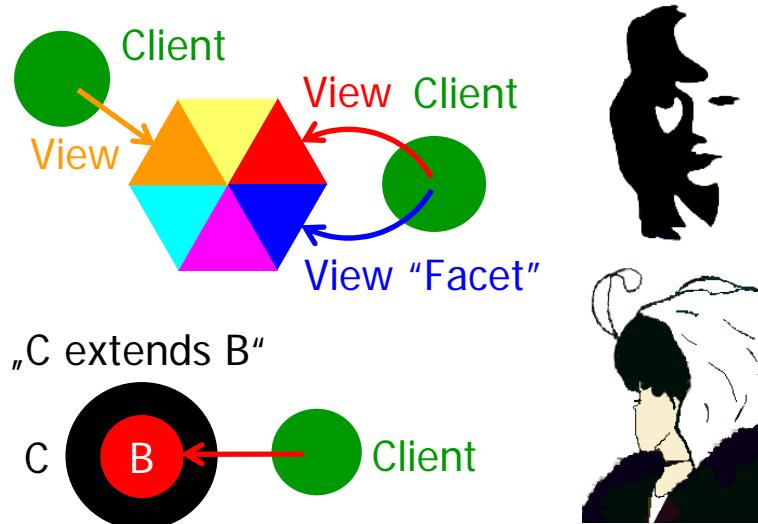
A Simple A&D Module

```
• module Zeller;
  var wd: array 7 of string;
  procedure { public } WeekDay
    (c, y, m, d: integer): string;
    var n: integer;
  begin
    n := entier(2.62*m - 0.2)
      + d + y + y div 4 + c div 4 - 2*c;
    return wd[n mod 7]
  end WeekDay;
begin
  wd[0] := "Sunday"; wd[1] := "Monday";
  wd[2] := "Tuesday"; wd[3] := "Wednesday";
  wd[4] := "Thursday"; wd[5] := "Friday";
  wd[6] := "Saturday"
end Zeller.
```

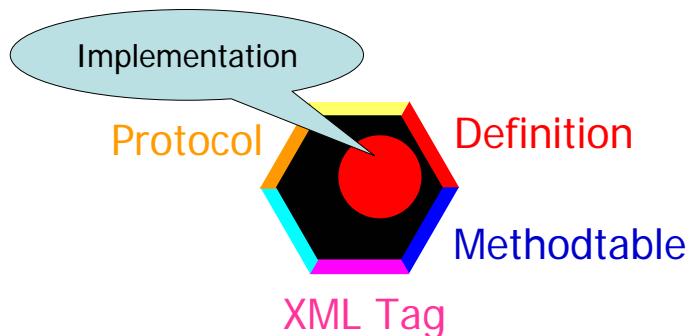
The Object Model

- Modular
- Compositional
- Active

Compositional vs. Hierarchical



Definition vs. Implementation



Example Jukebox (1)

```
definition Store;(*view*)
  procedure Clear;
  procedure Add (s: Songs.Song);
end Store.

implementation Store;
  var rep: Songs.Song;
  procedure Clear;
  begin rep := NIL
  end Clear;
  procedure Add (s: Songs.Song);
  begin s.next := rep; rep := s
  end Add;
begin Clear
end Store.
```

Example Jukebox (2)

```
definition Player;
  var cur: Songs.Song;
  procedure Play (s: Songs.Song);
  procedure Stop;
end Player.

object JukeBox implements Store, Player;
  (*aggregates implementation Store*)
  procedure Play (s: Songs.Song) implements
    Player.Play;
  begin ...
  end Play;
  procedure Stop implements Player.Stop;
  begin ...
  end Stop;
end JukeBox.
```

The Object Model

- Modular
- Compositional
- Active

The Challenge of Concurrency

- Moore's Law
 - Double performance each 1.5 years
- Achieved via
 - Until now: # FLOPS
 - 10 MHz → 100 MHz → 1 GHz → 3.2 GHz
 - Power, Heat ⇒ Stop at \approx 3.5 GHz
 - From now: Multi CPU cores
 - 1 CPU → 2 CPU → 8 CPU →
- Challenge of exploiting multiple CPU
 - Support needed from programming language

Some Language Design Goals ...

- Integrate concurrency with OOP
- Replace library calls with language constructs
- Abstract from deployment details (central or distributed)
- Present active objects as self-contained units with programming-language independent interfaces

Two New Constructs

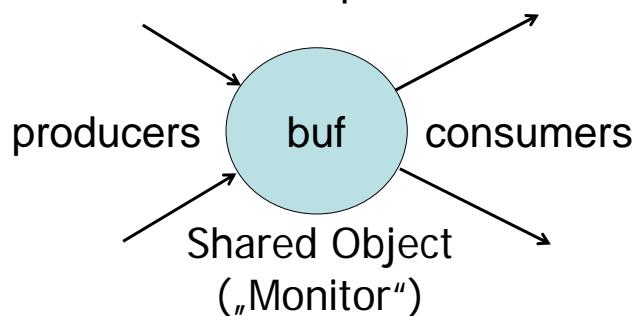
- The *await* statement
- Activities

The *await* Statement

- Used in shared objects for waiting on a local condition to be established by other activities
- Replaces the method of signalling by an autonomous concept

Example: Finite Buffer

- Scenario: consumers and producers communicating via finite buffer
- m = number of free slots in buffer
- n = number of occupied slots in buffer



Finite Buffer with Signals

- ```
public void Put (object x) {
 lock(this) {
 while (m == 0) { Monitor.Wait(this); }
 m--; buf[tail] = x;
 tail = (tail + 1) % size;
 n++; Monitor.PulseAll(this);
 }
}
```
- ```
public object Get () {
    lock(this) {
        while (n == 0) { Monitor.Wait(this); }
        n--; object x = buf[head];
        head = (head + 1) % size;
        m++; Monitor.PulseAll(this); return x;
    }
}
```

Cost: unnecessary context switches

Finite Buffer with *await*

- ```
procedure Put (var object x);
begin
 await (m # 0);
 dec(m); buf[tail] := x;
 tail := (tail + 1) mod size;
 inc(n);
end Put;
```
- ```
procedure Get (): object;
var x: object;
begin
    await (n # 0);
    dec(n); x := buf[head];
    head := (head + 1) mod size;
    inc(m); return x
end Get;
```

Activities

- Generalization of the procedural paradigm
- Used for multiple purposes
 - Run independent statements concurrently
 - Run intrinsic activities encapsulated in objects
 - Carriers of communications

Procedural Paradigm

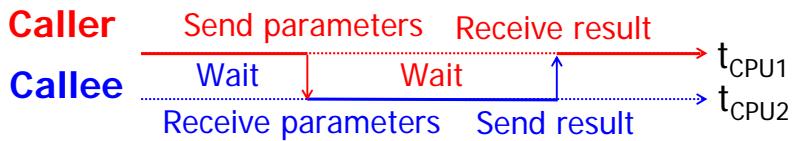
- Directed at single CPU configurations



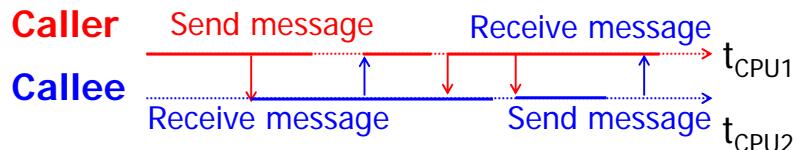
- The same paradigm used in different cases
 - Local procedure call
 - Method call
 - Remote procedure call

New Concept: Activities

- Procedure call as dialog



- Activities as generalized procedures



Scenario 1: Independent Actions

- **activity** A (...);
 var ...
 begin ...
 end A;
- **activity** B (...);
 var ...
 begin ...
 end B;
- **begin** { **barrier** }
 new A(...); new B(...)
end

Example 1: Quicksort

- **activity** Sort (l, h: integer)
 var i, j: integer;
 begin { **barrier** }
 ... (*partition array l, j & i, h*)
 if l < j **then new** Sort(l, j) **end**;
 if i < h **then new** Sort(i, h) **end**
 end Sort;
- (*start Quicksort*)
 begin
 new Sort(1, N)
 end

Example 2: Active Objects

- **type** X = **object**
 activity A (...);
 ... (*intrinsic behavior*)
 end A;
 activity B (...);
 ... (*intrinsic behavior*)
 end B;
 procedure new (...);
 ... (*constructor*)
 end new;
 begin { **barrier** }
 new A(...); **new** B(...)
 end X;

Passive vs. Active



- put item
- get item



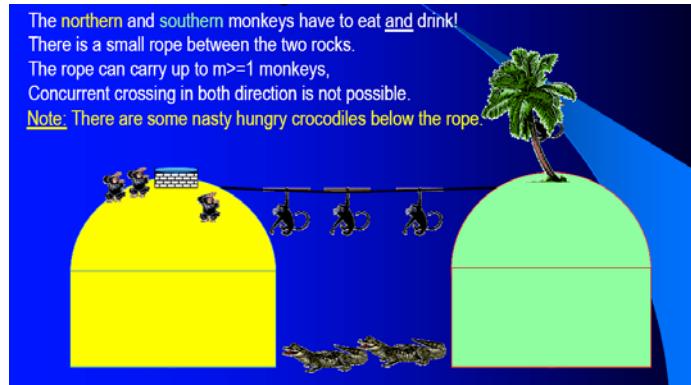
- set time
- get time

Scenario 2: Object Dialogs

```
• type Y = object (*callee*)
  activity D (...): ...;
  var t, u: T;
  begin (*dialog*)
    ... return t; ... u := *; ...
  end
end Y;

• var y: Y; d: Y.D; t, u: T;
begin (*caller*)
  y := new Y;
  d := new y.D; (*active link*)
  t := d(...); ... d(u); ...
end
```

Example 1: World of Monkeys



The Rope as Shared Resource

```
• module { shared } Rope; (*global view*)
  type
    Monkey = object; (*active*)
      MonkeyMsg = (claim, release);
    var cur, i: integer;
      (*number of monkeys on rope
       > 0 South-North traversal
       < 0 North-South traversal*)
    activity MonkeyDialog () : MonkeyMsg;
  begin
    for i := 0 to 99 do new Monkey () end
  end Rope;
```

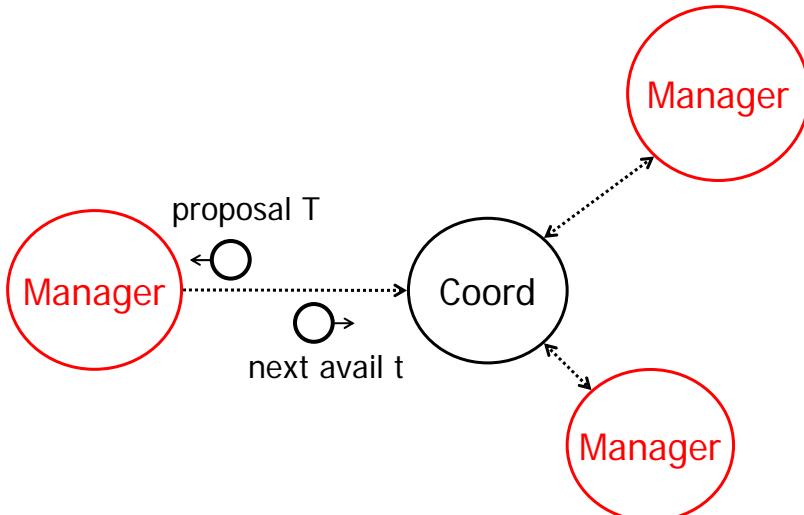
Monkeys as Active Objects

```
• type Monkey = object
  activity LiveOnTheRocks ();
    var res: MonkeyMsg;
        d: Rope.MonkeyDialog;
    begin (*story of life*)
      d := new Rope.MonkeyDialog;
      loop
        → passivate(Random.Next()); (*eat/dr*)
        res := d(MonkeyMsg.claim);
      end
    end LiveOnTheRocks;
  begin { barrier }
    new LiveOnTheRocks()
  end Monkey;
```

The Monkey Dialog Activity

```
• activity MonkeyDialog (): MonkeyMsg;
  var req: MonkeyMsg;
begin
  loop
    req := *; (*South-North request*)
    → await (0 <= cur) & (cur < m);
    inc(cur); passivate(100);
    dec(cur); return MonkeyMsg.release;
    req := *; (*North-South request*)
    await (0 >= cur) & (cur > -m);
    dec(cur); passivate(100);
    inc(cur); return MonkeyMsg.release
  end
end MonkeyDialog;
```

Example 2: Next Meeting Time



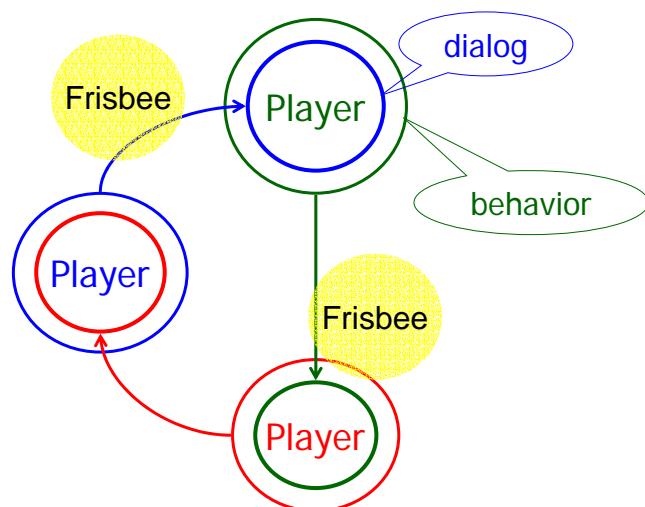
The Coordinator

```
• module { shared } Coordinator;
  type Manager = object; (*active*)
  var T, i: integer;
  activity ManagerDialog ();
  var next: integer;
begin
  loop
    return T; t := *;
    if t > T then T := t end;
    await T > t;
  end
  end ManagerDialog;
begin T := 0;
  for i := 0 to 9 do new Manager() end
end Coordinator.
```

Managers as Active Objects

```
• type Manager = object
  activity Check ();
  var t: integer;
    d: Coordinator.ManagerDialog;
begin
  d := new Coordinator.ManagerDialog;
loop
  t := d();
  (*check agenda and update t*)
  d(t)
end
begin new Check()
end Manager;
```

Example 3: Frisbee Fun



Starting the Game

```
• module Game;
  type Player = object; (*active*)
    var i: integer; p, q, last: Player;
begin
  last := new Player(); q := last;
  for i := 0 to 9 do
    p := new Player();
    p.Init(q, Random.Next() mod 2);
    q := p
  end;
  last.Init(q, 0)
end Game.
```

Player as Dual Activity Object

```
• type Player = object { shared }
  FrisbeeMsg = (request, catch);
  var nofFrisbees: integer;
      d: Player.FrisbeeDialog;
  procedure Init (q: Player; f: integer);
begin
  d = new q.FrisbeeDialog;
  nofFrisbees = f;
end Init;
activity Play ();
activity FrisbeeDialog (): FrisbeeMsg;
begin { barrier }
  new Play ()
end Player;
```

The Playing Activity

```
• activity Play ();
    var msg: FrisbeeMsg;
begin
    d := new Player.FrisbeeDialog;
    loop
        await nofFrisbees # 0;
        msg := d(); d(FrisbeeMsg.catch);
        nofFrisbees := 0
    end
end Play;
```

The Frisbee Dialog Activity

```
• activity FrisbeeDialog ();
    var msg: FrisbeeMsg;
begin
    loop
        await nofFrisbees = 0;
        return FrisbeeMsg.request;
        msg := *;
        nofFrisbees := 1
    end
end FrisbeeDialog;
```

Example 4: Santa Claus

- Invented by John Trono in „J. A. Trono. A new exercise in concurrency. SIGCSE Bulletin, 1994“
- Discussed and solved later by Ben-Ari with Rendez-Vous (in Ada95) and monitors (in Java)



The Original Story

- Santa Claus sleeps at the North pole until awakened by either all of the nine reindeer, or by a group of three out of ten elves. He performs one of two indivisible actions:
 - If awakened by the group of reindeer, Santa harnesses them to a sleigh, delivers toys, and finally unharnesses the reindeer who then go on vacation.
 - If awakened by a group of elves, Santa shows them into his office, consults with them on toy R&D, and finally shows them out so they can return to work constructing toys.
- A waiting group of reindeer must be served by Santa before a waiting group of elves. Since Santa's time is extremely valuable, marshalling the reindeer or elves into a group must not be done by Santa.

Semaphore Approach (1)

```
-- Consult
for All_Waiting_Elves loop
    V(Elf_Wait);
end loop;
for All_Elves loop
    V(Invite_In);
end loop;
Consult;
for All_Elves loop
    V>Show_Out);
end loop;

loop
    P(Santa);
    if All_Reindeer_Ready then
        -- Deliver
    else -- All_Elves_Ready
        -- Consult

-- Deliver
for All_Waiting_Reindeer loop
    V(Reindeer_Wait);
end loop;
for All_Reindeer loop
    V(Harness);
end loop;
Deliver_Toys;
for All_Reindeer loop
    V(Unharness);
end loop;
```

Semaphore Approach (2)

One for every
Reindeer and Elf
(correspondingly)

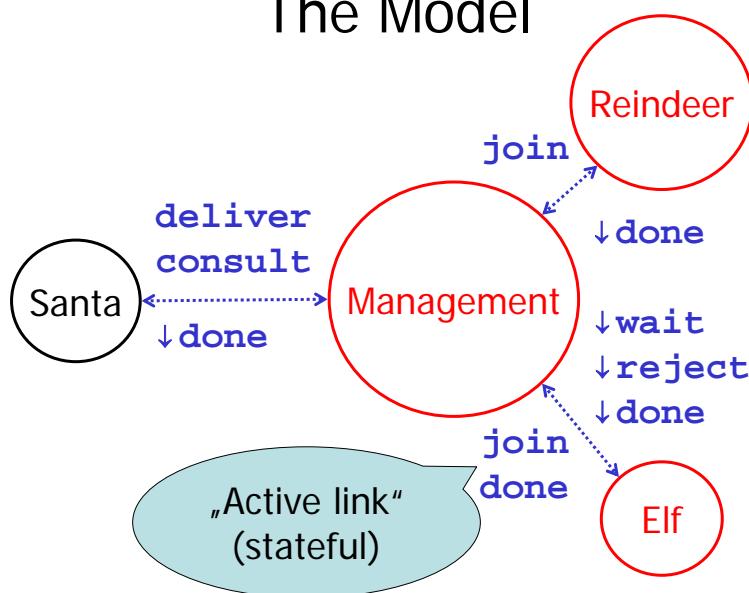
```
loop
    if Is_Last_Reindeer then
        V(Santa);
    else
        P(Reindeer_Wait);
    end if;
    P(Harness);
    Deliver_Toys;
    P(Unharness);
end loop;
```

Our Extension: Negotiation

- Before joining, elves should be informed about the expected waiting time and be given the opportunity to withdraw
- Dialog as formal syntax in EBNF
 - Messages from callee to caller marked "↓"

```
HandleElf = join ( Negotiate | ↓reject ).  
Negotiate = [ ↓wait join ] ↓done |  
↓wait done.
```

The Model



Management as Active Server

```
• module { shared } Management;
  type
    ElfMsg = (join, reject, delay, done);
    ReindeerMsg = (join, done);
    SantaMsg = (deliver, consult, done);
    Elf = object; (*active*)
    Reindeer = object; (*active*)
    Santa = object;
  var r0, r, R, e0, e, E: integer;
      santa: Santa;
  activity Work ();
  activity ElfDialog (): ElfMsg;
  activity ReindeerDialog (): ReindeerMsg;
begin { barrier }
  new Work ()
end Management.
```

Manager as Active Server

```
• activity Work ();
  var res: SantaMsg;
      d: Santa.Dialog;
begin
  d := new santa.Dialog;
  loop
    await (r > r0) & (e > e0);
    if r > r0 then
      res := d(SantaMsg.deliver); inc(r0)
    else
      res := d(SantaMsg.consult); inc(e0)
    end
  end
end Work;
```

The Elf Handling Activity

```
• activity ElfDialog ();
  var myGroup: integer; req: ElfMsg;
begin
  loop req := *;
    if (*too soon*) then return ElfMsg.reject
    else
      if e0 < e then
        return ElfMsg.wait; req := *
      end;
      if req = ElfMsg.join then
        myGroup = e; inc(E);
        if E = 3 then E := 0; inc(e) end;
        await e0 > myGroup;
        return ElfMsg.done
      end
    end
  end
end ElfDialog;
```

Elves as Active Objects

```
• type Elf = object
  activity Work ();
  var res: ElfMsg; d: Manager.ElfDialog;
begin
  d := new Manager.ElfDialog;
  loop
    passivate(Random.Next());
    res := d(ElfMsg.join);
    if res = ElfMsg.wait then
      if (*impatient*) then d(ElfMsg.done)
        else res := d(ElfMsg.join)
      end
    end
  end
end Work;
begin { barrier } new Work()
end Elf;
```

Santa Dialog Controlled

```
• type Santa = object
  activity Dialog (): SantaMsg;
  var req: SantaMsg;
begin
  loop
    req := *;
    if req = SantaMsg.deliver then
      passivate(10000)
    else (*consult*) passivate(500)
    end;
    return SantaMsg.done
  end
end Dialog;
end Santa.
```

Example 5: Rental System

```
module { shared } Rental;
const N = 100;
type
  Message = (accept, return);
  Client = object;
  var nofFree, i: integer;
    free: array N of boolean;
  activity Negotiate (): Message;
  procedure Next (obj: integer);
  procedure Release (obj: integer);
begin
  for i := 0 to N - 1 do free[i] := true end;
  nofFree := N
end Rental.
```

Rental System: Negotiate

```
activity Negotiate (): Message;
  var msg: Message; obj: integer;
begin
  obj := -1;
repeat
  obj = Next(obj); return obj;
  msg := *;
  if msg # Message.Accept then
    Release(obj)
  end
until msg = Message.Accept;
msg := * (*return*)
Release(obj)
end Negotiate;
```

Rental System: Find & Release

```
procedure Next (obj: integer);
begin
  await nofFree > 0;
  while ~free[obj] do
    obj := (obj + 1) mod N
  end;
  free[obj] := false; dec(nofFree);
  return obj;
end Next;

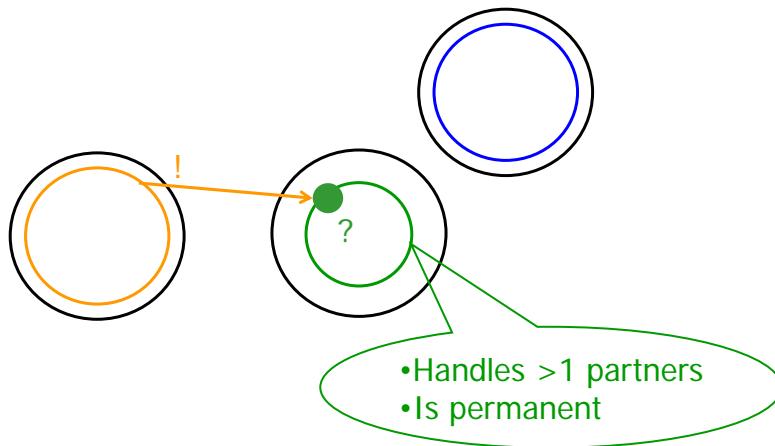
procedure Release (obj: integer);
begin
  free[obj] := true; inc(nofFree)
end Release;
```

Rental System: Client

```
type Client = object
  activity Rent ();
    var res: Message; d: Rental.Negotiate;
        suitable: boolean;
    begin d := new Rental.Negotiate;
    repeat
      obj := d(); suitable := Check(obj);
      if ~suitable then d(Message.Return) end
    until suitable;
    d(Message.Accept); (*now use the object*)
    d(Message.Return)
  end Rent;
  begin { barrier } Rent()
end Client;
```

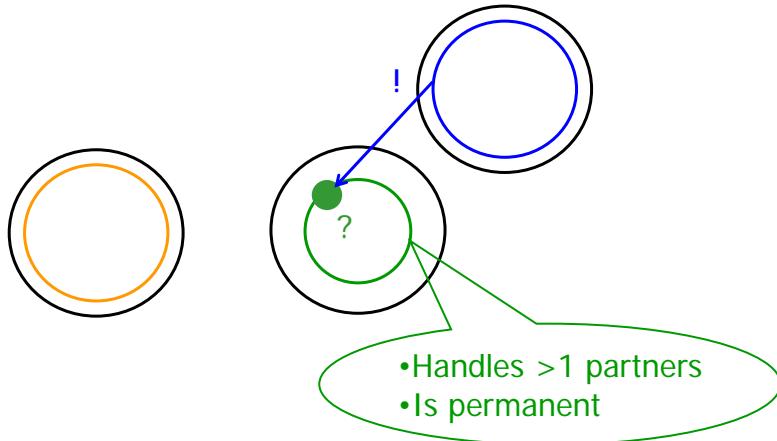
Active Links vs. CSP

- CSP Scenario 1



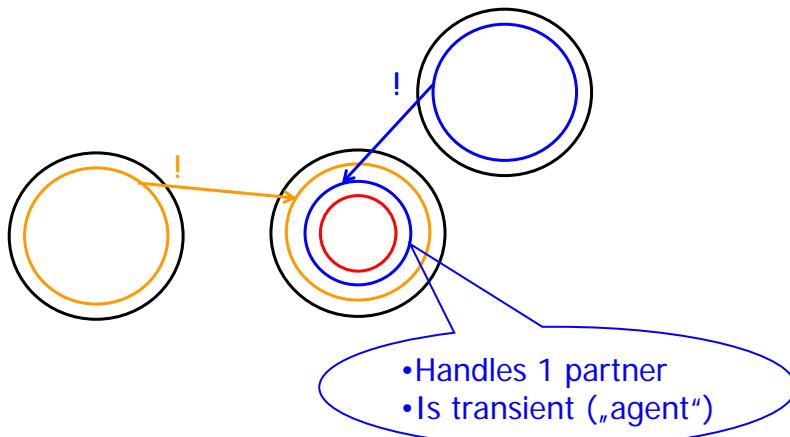
Active Links vs. CSP

- CSP Scenario 2



Active Links vs. CSP

- Active Links



Summary (1)

- The presented *await* construct
 - Adds autonomy to objects
 - Contributes to scalability
 - Delegates condition scheduling to the runtime system (compare with garbage collection)

Summary (2)

- The presented concept of activity upgrades the ordinary object-oriented model in three respects by adding
 - An option of orchestrating multiple concurrent activities according to programmed „launch logic“
 - Optional intrinsic encapsulated behavior of objects
 - A new way of dialog-oriented and stateful interoperability based on „active links“

Summary (3)

- The Zonnon concurrency model is an object-oriented combination of a shared-memory model and a message-passing model

Summary (4)

- The concept of activity has proved its suitability in several case studies and in implementations
 - The model of active objects underlies the *Aos Active Oberon* operating system
 - Active objects and dialogs have been implemented in the *Active C# ROTOR* compiler available from
<http://www.avocado.ethz.ch/ActiveCSharp/>
 - Activities are currently being implemented in the *Zonnon for .NET* compiler, see
<http://zonnon.ethz.ch>