

# Good Ideas - Revisited

Niklaus Wirth

September 2005

Moscow State University

Computer Architecture

Programming Languages

Miscellaneous Techniques

Programming Paradigms

# Expression stacks

$$(a/b) + ((c+d)*(c-d))$$

a b / c d + c d - \* +

a
b

a/b
-----

a/b
c
d

a/b
c+d

a/b
c+d
c
d

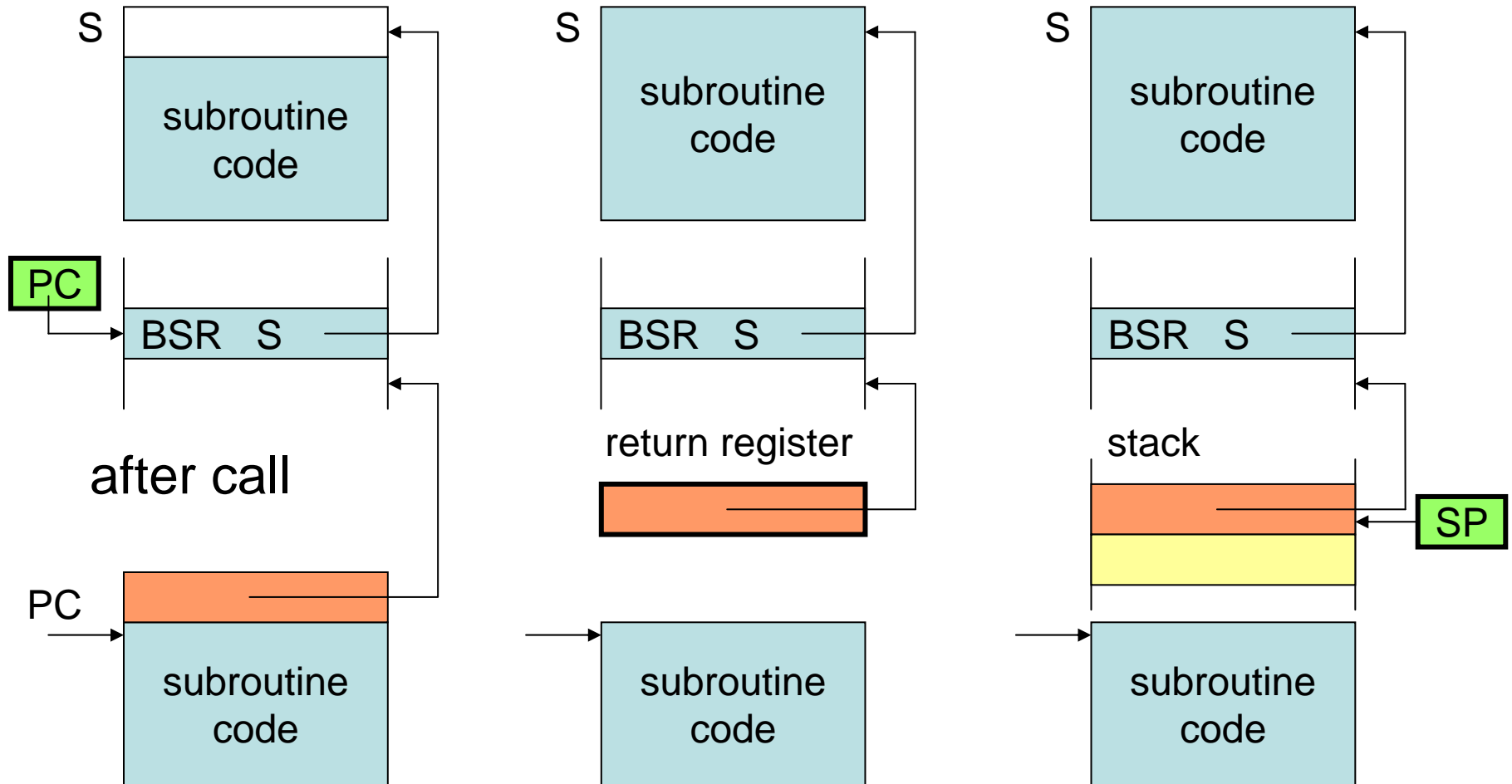
a/b
c+d
c-d

a/b
(c+d)*(c-d)

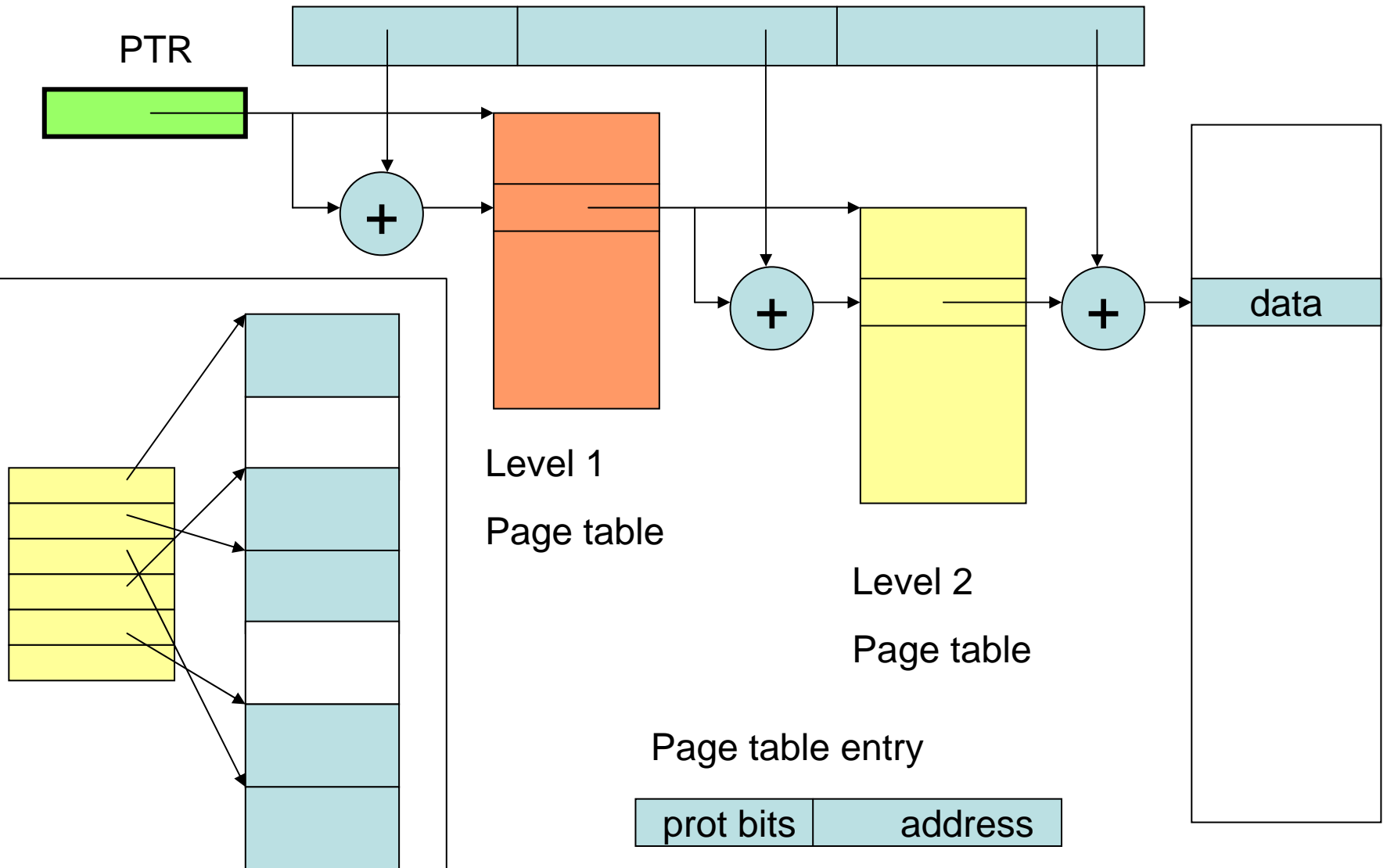
(a/b)+((c+d)*(c-d))
---------------------

# Subroutine return addresses

before call



# Virtual addressing



# Benefit and cost of virtual addressing

## Benefits:

- Simplifies allocation in multiprocessing
- Simplifies reuse of pages
- Protects from illegal access

## Cost:

- Multiple memory access via page tables
- Additional hardware
- Efficiency requires caches

# Simple vs. complex instruction sets

Simple instruction set -> simple hardware

What led to complex instruction sets?

- High-level language constructs

- Wish for code density

- Microprogramming

- Descriptor-architectures (B5500)

The concept of computer families

The return to RISC architectures

# Programming Language Features

- Notation and syntax
- Algol's FOR statement
- Algol's OWN variables
- Algol's name parameter

# Notation and Syntax

- Assignment operator  $x = y$     $x := y$

$z = x ++ y$     $b = x == y$

$x +++++y$     $++x++++y+1$

$x+++++y+1==++x++++y$

$x+++y++==x+++++y+1$

- APL:  $x-y-z$  means  $x-(y-z)$ , but not  $(x-y)-z$
- Confusion between statement and expression
- **Statements** are **executed**
- **Expressions** are **evaluated**



# Syntax (Algol and Pascal)

**if b then S0**                      **if b then S0 else S1**

**if b0 then if b1 then S0 else S1**

**if b0 then if b1 then S0 else S1**

**if b0 then if b1 then S0 else S1**

**if p then for i := 1 step 1 until n do if q then S1 else S2**

**if p then for i := 1 step 1 until n do if q then S1 else S2**

**if b then S0 end**                      **if b then S0 else S1 end**

# Algol's complicated **for** statement

for  $i := 1$  step 1 until  $n$  do  $a[i] := 0$

for  $i := 2, 3, 5, 7, 11$  do  $a[i] := 0$

for  $i := x, x+1, x*(y+z)$  do  $a[i] := 0$

for  $i := i+1$  while  $i < n$  do  $a[i] := 0$

for  $i := x-3, x$  step  $k$  until  $y, y+7$ , while  $z < 20$   
do  $a[i] := 0$

for  $i := 1$  step 1 until  $i+1$  do  $a[i] := 0$

for  $i := 1$  step  $i$  until  $i$  do  $i := -i$

# Algol's **own** variables

```
procedure P(x, y);  
begin integer z; z := x; x := y; y := z  
end
```

```
real procedure random;  
begin own real x; x := (x*a + b) mod c;  
  randon := x  
end
```

# Algol's name parameter

**real procedure** square(x); real x; square := x\*x

square(a) literally means a\*a

square(sin(a)\*cos(a)) stands for  
sin(a)\*cos(a)\*sin(a)\*cos(a)

**real procedure** square(x);

**value** x; real x; square := x\*x

stands for

**begin** real x'; x' := x; square := x'\*x' **end**

# Algol's Jensen device

```
real procedure sum(k, x, n);  
begin real s;  
  for k := 1 step 1 until n do s := s+x;  
  sum := x  
end
```

$a_1 + a_2 + \dots + a_{100}$

sum(i, a[i], 100)

$a \times b$

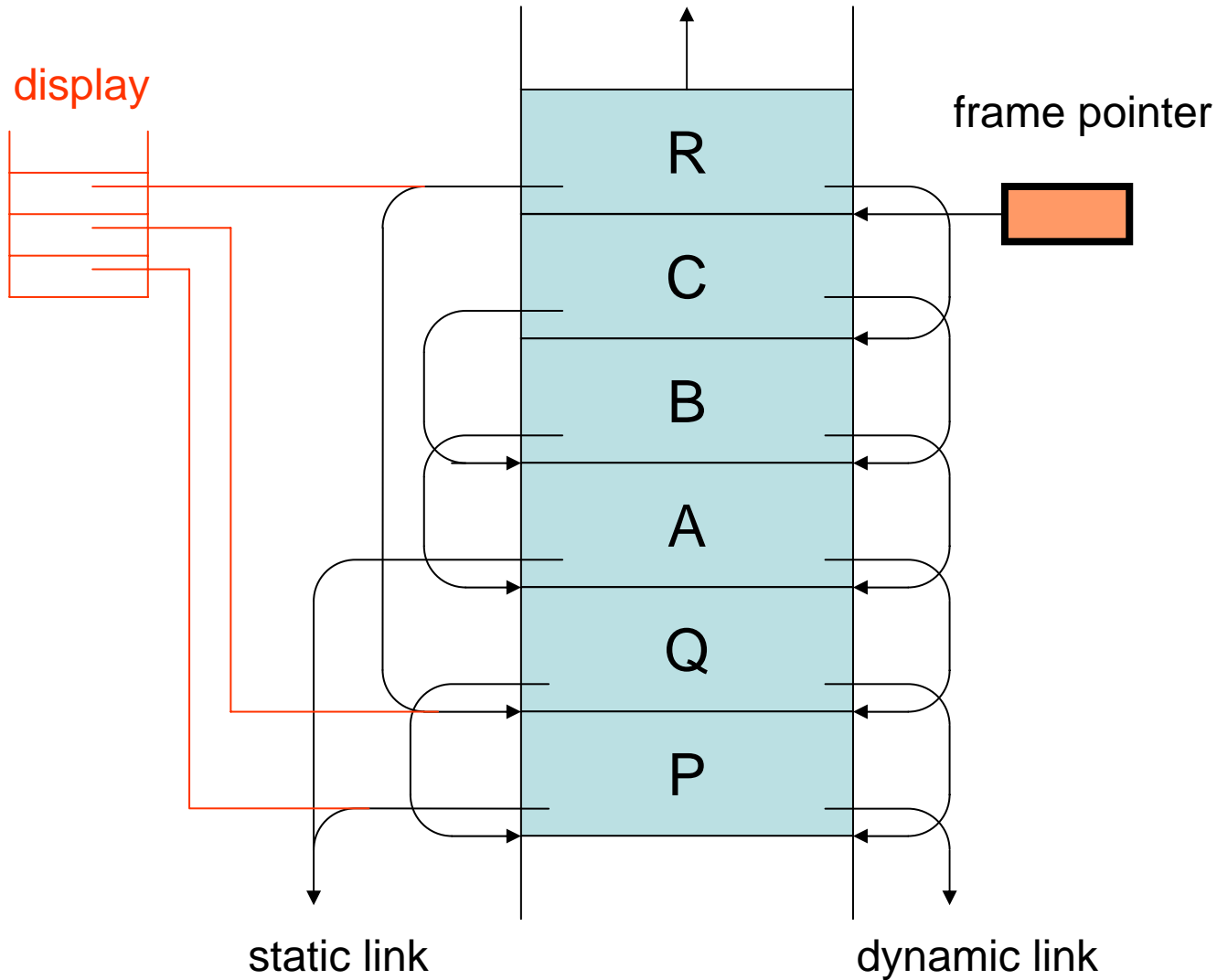
sum(i, a[i]\*b[i], 100)

# Dijkstra's "display"

```
procedure A(proc h);  
begin integer x;  
  procedure B;  
  begin integer y;  
    procedure C;  
    begin integer z; h  
    end ;  
    C  
  end ;  
  B  
end
```

```
procedure P;  
begin integer i;  
  procedure Q;  
  begin integer j;  
    procedure R;  
    begin integer k;  
    end ;  
    A(R)  
  end ;  
  Q  
end
```

→ P → Q → A → B → C → R



# Functional programming

- What distinguishes a functional language from a procedure language?
- FP: exclusively function application
- **No state, no variables**
- FPLs have sneaked in state and assignment
- **No side effects**
- Eases detection of potential parallelism
- Academic exercise ?



# Logic programming

- Prolog
- Implementation is a search engine for solutions satisfying given predicates
- Logic inference engine
- In practice requires hints in the form of cuts
- One must understand the functioning of the hidden engine
- An academic exercise ?

# Object-oriented programming

- Technically based on 2 concepts only:
  - procedure types of variables
  - type extension (inheritance)
- View of procedures as belonging to objects
- Terminology

**object** (record typed) variable

**class** (record) type

**method** (record bound) procedure

**send msg** call procedure

# “Good ideas” of today?

- Many of the good ideas of their time have become mediocre or even bad ideas
  - Because of technology changes
  - Because of shifts of goals and habits
  - Because of too much emphasis on efficiency
- Which are today’s “good ideas”?
  - Will they also turn mediocre?
  - Or are they bad already now?